

Course Name: A Level (1<sup>st</sup> Sem)

Subject : Introduction to DBMS

Topic: DB Normalization – Exercise Practice on 2NF (Part 5)

Date: 19-May-2020

### Database Normalization – Exercise practices on 2NF

#### Que 1:

Suppose a following stu\_proj relational schema:

| stu_id | proj_id | stu_name | proj_name       | proj_lang_used |
|--------|---------|----------|-----------------|----------------|
| s101   | p001    | Rakesh   | Online Chatting | python         |
| s102   | p001    | Kritika  | Online Chatting | python         |
| s102   | p002    | Kritika  | Text Editor     | Java           |
| s103   | p002    | Mahesh   | Text Editor     | Java           |
| s104   | p002    | Ram      | Text Editor     | Java           |
| s104   | p003    | Ram      | Online Shopping | PHP            |

Identify functional dependencies in above table and find out that relation is in 2NF or not? If not decompose it in 2NF.

(A student can work on many projects and a project can have many students associated with it.)

#### Solution:

The table is in 1NF because each attribute in the table have atomic (single) value.

The following FDs are identified based on the value and descriptions given about table:

**stu\_id** → **stu\_name** ( Student name can be determined by student id)

**proj\_id** → **proj\_name, proj\_lang\_used** (Project name and language used can be determined by project id)

Since (stu\_id, proj\_id) uniquely identifies each record in the table, and the closure of (stu\_id, proj\_id) also has all the attributes of table.

So, (stu\_id, proj\_id) is the candidate key in table.

**Prime attributes** - stu\_id, proj\_id ( because these are the part of the candidate key)

**Non prime attributes** - stu\_name, proj\_name, proj\_lang\_used (because these are not part of the candidate key)

Now, the partial dependency exists in the table, it means violation of rule of 2NF because non prime attribute is dependent on part of the candidate key.

|                                     |   |
|-------------------------------------|---|
| stu_id → stu_name                   | both FD violating rule of 2NF, since (stu_id, proj_id) is the candidate key |
| proj_id → proj_name, proj_lang_used |   |

**Therefore**, to convert the relation in 2NF, It is divided into three relations:

|   |   |
|---|---|
| <b>student</b> (stu_id, stu_name)                   | Since stu_id → stu_name                   |
| <b>project</b> (proj_id, proj_name, proj_lang_used) | Since proj_id → proj_name, proj_lang_used |
| <b>stu_proj_alloc</b> (stu_id, proj_id)             | Since stu_id, proj_id is candidate key    |

**Above three tables (student, project, stu\_proj\_alloc) are following the rules of 2NF.**

See, how the data redundancy is minimized after converting table into 2NF (compare it with original table):

table1: **student**

| stu_id | stu_name |
|--------|----------|
| s101   | Rakesh   |
| s102   | Kritika  |
| s103   | Mahesh   |
| s104   | Ram      |

table2: **project**

| proj_id | proj_name       | proj_lang_used |
|---------|-----------------|----------------|
| p001    | Online Chatting | python         |
| p002    | Text Editor     | Java           |
| p003    | Online Shopping | PHP            |

table3: **stu\_proj\_alloc**

| stu_id | proj_id |
|--------|---------|
| s101   | p001    |
| s102   | p001    |
| s102   | p002    |
| s103   | p002    |
| s104   | p002    |
| s104   | p003    |

**Exercise:**

1. Suppose a relational schema R (A B C), and

FDs:  $AB \rightarrow C$

$B \rightarrow C$

Check out the relation R is in 2NF or not? If not decompose it in 2NF.

