# National Institute of Electronics & Information Technology (NIELIT), Gorakhpur राष्ट्रीय इलेक्ट्रॉनिकी एवं सूचना प्रौद्योगिकी संस्थान ,गोरखपुर



Course Name: A Level (1<sup>st</sup> Sem)

Topic: **DB Normalization – Dependency Preserving** 

**Decomposition (Part 13)** 

Subject: Introduction to DBMS

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## **Database Normalization – Dependency Preserving Decomposition**

Suppose R is a relational schema and F is the set of functional dependencies on R.

If R is decomposed into relations R1, R2,  $\dots$  Rn , each holding functional dependencies F1, F2,  $\dots$  Fn respectively.

We can say,  $F' = F1 \cup F2 \cup ... \cup Fn$ 

Now this decomposition will be considered as **dependency preserving decomposition** if and only if

# Every dependency in F is logically implied by $F^*$ i.e. $F^{*+} = F^{+}$

It is obvious that  $F1 \subseteq F^+$ ,  $F2 \subseteq F^+$ , and so on

If we verify that F is satisfied in R, we have verified that decomposition is dependency preserving decomposition i.e.  $F1 \cup F2 = F$ 

Suppose R (A B C D) and set of functional dependencies

$$F: AB \rightarrow CD$$

$$D \rightarrow A$$

If R is decomposed into following two relations

R1 (A D),

R2 (BCD)

Identify that this decomposition is dependency preserving or not?



#### **Solution:**

$$R (A B C D)$$

$$F: AB \longrightarrow CD$$

$$D \longrightarrow A$$

Since there are two attributes in R1 i.e. A, D

so find A<sup>+</sup>, D<sup>+</sup>

[Closure must be calculated using FD set in R]

$$A^{+} = A$$
 [It concludes trivial FD A  $\rightarrow$  A]  
 $D^{+} = AD$  [It concludes FD D  $\rightarrow$  A]

The closure of attributes concludes, following FD exist in R1

$$F1 \{ D \rightarrow A \}$$

### R2 (B C D)

Since there are three attributes in R2 i.e. B, C, D

so find B<sup>+</sup>,C<sup>+</sup>, D<sup>+</sup>, BC<sup>+</sup>, CD<sup>+</sup>, BD<sup>+</sup>

[Closure must be calculated using FD set in R]

$$B^+=B$$
 [trivial]

$$C^+ = C$$
 [trivial]

$$BC^+=BC$$
 [trivial]

R2]

$$BD^+=BDAC$$
 [It concludes FD, BD  $\rightarrow$  C

A is not included at right side of FD because A is not

attribute in R2]

The closure of attributes concludes, following FD exist in R2

$$F2 \{ BD \rightarrow C \}$$

Now 
$$F' = F1 \cup F2 = \{ D \rightarrow A, BD \rightarrow C \}$$
  
and original  $F$  in  $R = \{AB \rightarrow CD, D \rightarrow A\}$ 

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It can be seen that while decomposing, the functional dependency AB → CD has been lost.

To understand, calculate  $AB^+ = AB$  [ Using FD in F`] while in F, AB can determine CD but this dependency is lost in F`.

Therefore, this decomposition is not dependency preserving decomposition.

### **Exercise:**

Suppose R (A B C D) and set of functional dependencies

$$F: AB \rightarrow C$$

$$C \rightarrow D$$

$$D \rightarrow A$$

If R is decomposed into following two relations

Identify that this decomposition is dependency preserving or not?

